## seL4: Experiences, Improvements, and Optimizations

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## Common seL4 Use Cases

### Isolate legacy systems in Virtual Machines •





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- Isolate legacy systems in Virtual Machines •
- **Cross-domain Solution** lacksquare





General Purpose VM

## Common seL4 Use Cases

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- **Cross-domain Solution**







Software implementation of a network device

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- Leverage existing network stacks to communicate between VMs •

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### How does it work? From the guest's perspective

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- Guest scans PCI bus and loads virtio-net driver

### with virtio-net device present io-net driver



### How does it work? From the guest's perspective

- Host implements a virtual PCI bus with virtio-net device present •
- Guest scans PCI bus and loads virtio-net driver

root@xilinx-zcu102-2021\_1:~# ip a 1: lo: <LOOPBACK,UP,LOWER\_UP> mtu 65536 qdisc noqueue qlen 1000 link/loopback 00:00:00:00:00:00 brd 00:00:00:00:00:00 inet 127.0.0.1/8 scope host lo valid\_lft forever preferred\_lft forever inet6 ::1/128 scope host valid\_lft forever preferred\_lft forever 2: eth0: <BROADCAST,MULTICAST,UP,LOWER\_UP> mtu 15964 qdisc pfifo\_fast qlen 1000 link/ether 00:00:00:00:00:01 brd ff:ff:ff:ff:ff:ff inet 192.168.1.1/24 scope global eth0 valid\_lft forever preferred\_lft forever inet6 fe80::200:ff:fe00:1/64 scope link valid\_lft forever preferred\_lft forever 3: sit0@NONE: <NOARP> mtu 1480 qdisc noop qlen 1000 link/sit 0.0.0.0 brd 0.0.0.0

Standard network device available for use by any networking application

Host implements a virtual PCI bus with virtio-net device present

- Host implements a virtual PCI bus with virtio-net device present •
- Guest accesses the PCI bus when transmitting a packet  $\bullet$





- Host implements a virtual PCI bus with virtio-net device present
- Guest accesses the PCI bus when transmitting a packet
- Host knows where the network packet is stored in memory



### Host receives a notification with a virtio-net specific badge

0

VirtioNet • now

Notification

You've got a packet in your virtqueue!



Se /

~

Security. Performance. Proof.

- Host receives a notification with a virtio-net specific badge •
- Host reads the packet from the virtqueues and into guest memory •



- Host receives a notification with a virtio-net specific badge
- Host reads the packet from the virtqueues and into guest memory •
- Host injects an interrupt to the guest



### How well does virtio-net work on seL4?

- Works great for simple communication path •
  - pings, status updates, etc



### How well does virtio-net work on seL4?

- Works great for simple communication path •
  - pings, status updates, etc
- Works poorly for throughput focused application

ro	oot@z	zcu102-zynqmp:	~# ip	erf3 -c 192.1	.68.1.2			
	onneo 51	cting to host local 192 168	192.1 1 1	68.1.2, port	5201 prected to 192 1	68 1 2	nort	5201
ľ	ID]	Interval		Transfer	Bitrate	Retr	Cwnd	5201
]	5]	0.00-1.00	sec	6.46 MBytes	54.2 Mbits/sec	14	60.8	KBytes
Ι	5]	1.00-2.00	sec	5.53 MBytes	46.4 Mbits/sec	26	43.8	KBytes
Ι	5]	2.00-3.00	sec	6.46 MBytes	54.2 Mbits/sec	5	35.4	KBytes
I	5]	3.00-4.00	sec	5.78 MBytes	48.5 Mbits/sec	13	39.6	KBytes
1	5]	4.00-5.00	sec	5.47 MBytes	45.9 Mbits/sec	28	7.07	KBytes
1	5]	5.00-6.00	sec	5.72 MBytes	48.0 Mbits/sec	29	31.1	KBytes
]	5]	6.00-7.00	sec	5.22 MBytes	43.8 Mbits/sec	37	14.1	KBytes
Ι	5]	7.00-8.00	sec	5.90 MBytes	49.5 Mbits/sec	13	32.5	KBytes
Ι	5]	8.00-9.00	sec	5.90 MBytes	49.5 Mbits/sec	14	65.0	KBytes
]	5]	9.00-10.00	sec	5.59 MBytes	46.9 Mbits/sec	24	8.48	KBytes
Ę.								
Ļ	ID	Interval		Transfer	Bitrate	Retr		
L	5]	0.00-10.00	sec	58.0 MBytes	48.7 Mbits/sec	203		sender
	5]	0.00-10.01	sec	57.4 MBytes	48.1 Mbits/sec			receive

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- Upstream methodology:
  - 1. Find physical address to read from
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  - 3. Use capability to map that page into Host's address space
  - 4. Read/write information using that page
  - 5. Unmap the page from the Host's address space

- How often does the host need to access the guest's memory?
  - A lot!
- 4 functions:
  - 1. ring\_avail reads guest memory once

```
uint16_t ring_avail(virtio_emul_t *emul, struct vring *vring, uint16_t idx)
{
    uint16_t elem;
    vm_guest_read_mem(emul->vm, &elem, (uintptr_t) & (vring->avail->ring[idx % vring->num]), sizeof(elem));
    return elem;
}
```



- How often does the host need to access the guest's memory?
  - A lot! •
- 4 functions:
  - 1. ring\_avail reads guest memory once
  - ring\_avail\_idx reads guest memory 2. once

```
uint16_t ring_avail_idx(virtio_emul_t *emul, struct vring *vring)
   uint16_t idx;
   vm_guest_read_mem(emul->vm, &idx, (uintptr_t)&vring->avail->idx, sizeof(vring->avail->idx));
   return idx;
```



- How often does the host need to access the guest's memory?
  - A lot! •
- 4 functions:  $\bullet$ 
  - 1. ring avail reads guest memory once
  - 2. ring avail idx reads guest memory once
  - 3. ring desc reads guest memory once

```
struct vring_desc ring_desc(virtio_emul_t *emul, struct vring *vring, uint16_t idx)
{
    struct vring_desc desc;
    vm_guest_read_mem(emul->vm, &desc, (uintptr_t) & (vring->desc[idx]), sizeof(desc));
    return desc;
```

- How often does the host need to access the guest's memory?
  - A *lot*!
- 4 functions:
  - 1. ring\_avail reads guest memory once
  - 2. ring\_avail\_idx reads guest memory once
  - 3. ring\_desc reads guest memory once
  - 4. ring\_used\_add reads guest memory
     once, writes guest memory twice

```
void ring_used_add(virtio_emul_t *emul, struct vring *vring, struct vring_used_elem elem)
{
    uint16_t guest_idx;
    vm_guest_read_mem(emul->vm, &guest_idx, (uintptr_t)&vring->used->idx, sizeof(vring->used->idx));
    vm_guest_write_mem(emul->vm, &elem, (uintptr_t)&vring->used->ring[guest_idx % vring->num], sizeof(elem));
    guest_idx++;
    vm_guest_write_mem(emul->vm, &guest_idx, (uintptr_t)&vring->used->idx, sizeof(vring->used->idx));
```



### **Guest Memory Access - Transmit Path**

- emul notify tx:
  - 1x: ring avail idx, ring avail, ring desc
  - 1x: direct read of guest memory to get network packet
- emul tx complete:
  - 1x: ring used add
- Total:
  - 7 maps, 7 unmaps

## Guest Memory Access - Receive Path

- emul rx complete:
  - 1x: ring avail idx, ring avail, ring desc
  - 1x: direct write of guest memory to give network packet
- Total:
  - 7 maps, 7 unmaps

## Can we speed this up?

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# How does this help?

### **Translation Vspace** Results

Throughput improves by around 8x!

[ 5]	local 192.168	.1.2	port 45814 (	connected to 192.	168.1.1	port	5201
[ ID]	Interval		Transfer	Bitrate	Retr	Cwnd	
[ 5]	0.00-1.01	sec	46.2 MBytes	s 385 Mbits/sec	0	199	KBytes
[ 5]	1.01-2.01	sec	45.0 MBytes	s 378 Mbits/sec	10	132	KBytes
[ 5]	2.01-3.02	sec	46.2 MBytes	s 382 Mbits/sec	0	164	KBytes
[ 5]	3.02-4.01	sec	45.0 MBytes	s 381 Mbits/sec	1	156	KBytes
[ 5]	4.01-5.01	sec	45.0 MBytes	s 381 Mbits/sec	4	150	KBytes
[ 5]	5.01-6.02	sec	46.2 MBytes	s 381 Mbits/sec	1	147	KBytes
[ 5]	6.02-7.02	sec	45.0 MBytes	s 381 Mbits/sec	0	189	KBytes
[ 5]	7.02-8.02	sec	45.0 MBytes	s 375 Mbits/sec	5	163	KBytes
[ 5]	8.02-9.01	sec	45.0 MBytes	s 381 Mbits/sec	0	195	KBytes
[ 5]	9.01-10.01	sec	45.0 MBytes	s 380 Mbits/sec	0	198	KBytes
[ ID]	Interval		Transfer	Bitrate	Retr		
[5]	0.00-10.01	sec	454 MBytes	s 380 Mbits/sec	21		sende
[5]	0.00-10.01	sec	454 MBytes	s 380 Mbits/sec			recei



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[ 5]	3.02-4.01	sec	45.0 MBytes	s 381 Mbits/sec	1	156	KBytes
[ 5]	4.01-5.01	sec	45.0 MBytes	s 381 Mbits/sec	4	150	KBytes
[ 5]	5.01-6.02	sec	46.2 MBytes	s 381 Mbits/sec	1	147	KBytes
[ 5]	6.02-7.02	sec	45.0 MBytes	s 381 Mbits/sec	0	189	KBytes
[ 5]	7.02-8.02	sec	45.0 MBytes	s 375 Mbits/sec	5	163	KBytes
[ 5]	8.02-9.01	sec	45.0 MBytes	s 381 Mbits/sec	0	195	KBytes
[ 5]	9.01-10.01	sec	45.0 MBytes	s 380 Mbits/sec	0	198	KBytes
[ ID]	Interval		Transfer	Bitrate	Retr		
[5]	0.00-10.01	sec	454 MBytes	s 380 Mbits/sec	21		sende
[5]	0.00-10.01	sec	454 MBytes	s 380 Mbits/sec			recei



### With Great Throughput Comes... More bugs

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  - More data going through virtqueues -> increase virtqueue size, etc.

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### With Great Throughput Comes... More bugs

- aren't seen at lower speeds
- Most were simple to fix:
  - More data going through virtqueues -> increase virtqueue size, etc.
- Others, not so much:  $\bullet$ 
  - corruption

### Increasing the throughput introduced a number of bugs into the system that

Increased throughput clobbers the cache and can lead to descriptor



## **Optimizing the Virtio-Net Driver**





## **Optimizing the Virtio-Net Driver**





**—** 

### Virtqueues

Module Layer

Virtqueues

### **Optimized Memcpy** Results

Connecting to host 192.168.1.100, port 5201									
[ 5]	local 192.168	3.1.10	1 port 36190	connected to 192	2.168.1	.100 p	oort 5201		
[ ID]	Interval		Transfer	Bitrate	Retr	Cwnd			
[ 5]	0.00-1.01	sec	64.9 MBytes	541 Mbits/sec	0	260	KBytes		
[ 5]	1.01-2.02	sec	62.4 MBytes	518 Mbits/sec	0	273	KBytes		
[ 5]	2.02-3.01	sec	61.2 MBytes	521 Mbits/sec	0	287	KBytes		
[ 5]	3.01-4.02	sec	62.5 MBytes	519 Mbits/sec	0	287	KBytes		
[ 5]	4.02-5.00	sec	61.2 MBytes	519 Mbits/sec	0	287	KBytes		
[ 5]	5.00-6.01	sec	62.5 MBytes	520 Mbits/sec	0	287	KBytes		
[ 5]	6.01-7.02	sec	62.5 MBytes	521 Mbits/sec	0	287	KBytes		
[ 5]	7.02-8.01	sec	61.2 MBytes	520 Mbits/sec	0	287	KBytes		
[ 5]	8.01-9.01	sec	62.5 MBytes	520 Mbits/sec	0	287	KBytes		
[ 5]	9.01-10.00	sec	61.2 MBytes	520 Mbits/sec	0	287	KBytes		
[ ID]	Interval		Transfer	Bitrate	Retr				
[ 5]	0.00-10.00	sec	622 MBytes	522 Mbits/sec	0		sender		
[ 5]	0.00-10.01	sec	622 MBytes	521 Mbits/sec			receive		





## What's next?

- **DornerWorks is open-sourcing the Translation Vspace and memcpy** ulletoptimizations
- We also developed a further improvement available for purchase...

### What's next?

Throughput (Mbps)

root@	xilinx-zcu102-	-2021_	1:~# iperf3 -	c 192.168.1.2			
Conneo	cting to host	192.1	68.1.2, port	5201	명이 있었다. 가 다 이 말 같은 것이다.	일 한 가지 않는 것 같은 것 같	
[5]	local 192.168	3.1.1	port 58794 co	nnected to 192.1	L68.1.2	port 5201	
[ ID]	Interval		Transfer	Bitrate	Retr	Cwnd	
[5]	0.00-1.01	sec	179 MBytes	1.49 Gbits/sec	6	357 KBytes	
[5]	1.01-2.00	sec	172 MBytes	1.46 Gbits/sec	19	202 KBytes	
[5]	2.00-3.00	sec	176 MBytes	1.48 Gbits/sec	8	295 KBytes	
[5]	3.00-4.01	sec	176 MBytes	1.47 Gbits/sec	2	280 KBytes	
[5]	4.01-5.00	sec	176 MBytes	1.49 Gbits/sec	11	264 KBytes	
[5]	5.00-6.00	sec	176 MBytes	1.47 Gbits/sec	4	373 KBytes	
[5]	6.00-7.00	sec	175 MBytes	1.47 Gbits/sec	6	295 KBytes	
[5]	7.00-8.01	sec	176 MBytes	1.47 Gbits/sec	4	280 KBytes	
[5]	8.01-9.00	sec	175 MBytes	1.47 Gbits/sec	4	295 KBytes	
[5]	9.00-10.01	sec	176 MBytes	1.47 Gbits/sec	4	280 KBytes	
				·			
[ ID]	Interval		Transfer	Bitrate	Retr		
[5]	0.00-10.01	sec	1.72 GBytes	1.47 Gbits/sec	68	send	er
[ 5]	0.00-10.02	sec	1.72 GBytes	1.47 Gbits/sec		rece	iver





